

Energy Efficient Routing by Reducing Iterative Broadcasting Mechanism in MANETs

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Abstract-Mobile ad hoc network composes of independent devices communicating with each other directly without aid of any centralized administration. In general some nodes in network stop working because they run out of battery power, therefore it is much important to preserve the power. In MANETs the nodes are portable so routing and power managing has turn into serious issue. Nodes in MANETs have limited battery capacity. In this work, we recommend an energy-efficient routing protocol for MANETs to decrease energy depletion during Route Discovery process. Previous mechanisms mainly engrossed on the shortest path schemes to reduce energy, which might end in network failure. This can lead to energy discrepancy and network life reduction. We propose an energy efficient routing protocol which increase network lifetime and consumes less power.

Keywords: MANET; Route Discovery ; Battery capacity; Energy Efficiency; Network Lifetime; Energy Consumption

1. INTRODUCTION

Due to MANETs prospective uses in various circumstances such as a military battle field, disaster relief, , etc., Mobile Ad hoc networks [1, 2] have newly arose as a primary research issue. MANETs consist of hosts that interconnect without a fixed structure. Communications takes place within the transmission range over a wireless channel, where each host connects with others in the neighborhood. Due to restricted infrastructure, dynamic connectivity and multi hop level relay based communication, achieving quality of service in an ad hoc network is a challenging subject. The lifespan of the network is a practical challenge in MANETs as the battery level of the nodes are delimited to fixed level. Network is partitioned if the energy level of the battery is zero. Hence the energy preservation methods are important to improve lifetime of the battery. It is necessary that energy proficient routing protocols should be used. All the conventional transmitting approaches fails to look after the wanted proportionality between battery levels required and available and uses shortest path routing. If the paths are used very high in the network, the nodes in that particular path die because of their battery capacity. In this regard, power aware routing protocols have been designed which select the routes to reserve battery lifespan of the nodes. The power aware routing approaches select paths such a way that packet routes through energy efficient path. The ideal proposal of energy efficiency under these circumstances is a vital requirement for MANET and focuses on the most cost-effective way to use battery power while guaranteeing appropriate operation of the network. The fundamental method of route discovery and best path selection are taken from routing protocols such as DSR [2], TORA [3], and AODV [4].

2. RELATED WORKS

In shortest path routing the power depletion at each hop nodes is the issue that fascinates present research. In this respect, many power aware routing elucidations are projected. Projected an innovative means to minimize power use in Shu-Lin Wu et al [5] spite of intensifying channel use. To measure the battery lifespan this scheme considers the RTS/CTS packet transmissions. S. Singhet al [6] proposed the PAMAS protocol, which targets to avoid the depletion of the lifespan of battery of a node during idle time. This protocol used two distinct channels one is for information and other is for signal. The indicator channel alerts the nodes to halt their RF devices during idle time. Further, Wan et al., [7] and Govardhan et al [17] suggested a least power usage routing which is dedicated to fixed MANETs. Our planned work concerned with power-conscious method for MANET routing protocols.

Theodore S et al [10], proposed the power is directly proportional to d^α , here ' d ' signifies the distance between two nodes and ' α ' value is in the range of 2 to 4. So MTPR route discovery process selects more nodes with less distances. But the start to end delay is more when compared to the power usage. In this, the chosen route frequently experience the splintering due to nodes with no battery level. Li and Wan [9] described a protocol which constructs a topology of least power and developed an process to find a shortest path whose distance is within a constant factor of that path. Energy consumption is calculated in terms of the length of the path. This proposed algorithm used only local information

3. Energy Efficient Routing by Reducing Iterative Broadcasting Mechanism (EER-RBM)

In the projected EER-RBM methodology, each host in the network has a routing table which has the capacity of battery of its neighbor nodes, route request status during RT discovery. At this point we suggest a new restricted dissemination to decrease overhead of packet that picks best paths with least distance and minimum energy. We consider a new method that confesses the hosts to exchange their existing battery level with their neighbor.

3.1. RT Discovery Phase

A node selects the neighbour to broadcast a RT by taking into account the ratio between life of battery required and life of battery available at hop level neighbor nodes.

3.2. RT Request (RReq)

The source host ' s ' routes the message by measuring the battery level ' bc ' which is calculated as the average energy indulgence per bit of data that transfers to its hop level nodes. Along with the RT request (RReq) the source node transmits that ' bc '. When each node that gets a RT request in a dissemination way it validates the status of the RT request. If that RT request is received previously then it rejects otherwise it verifies, whether that node itself is destination node, else it begins conditional broadcasting of that RT request to its hop level nodes.

The route request RReq structure is as follows:

$$RReq = \langle id(s), id(d), bc, \{h_1, h_2, \dots, h_i\} \rangle$$

3.3. RT Reply (RRep)

Upon getting the RT request RReq, the destination host ' d ' verifies RReq status. If it is replicate of the earlier received RReq it discards otherwise, the destination host ' d ' sends the RRep. The RRep that is initiated, uses the same path that is used by RReq but in the reverse direction. Each hop level node h_i updates the RRep with the distance $d(h_i \rightarrow h_{i+1})$ upon receiving the RRep. This process is repeated until RRep is received by the source node ' s '.

The RRep configuration of a hop level node h_i is as follows:

$$RRep = \langle id(s), id(d), hl, dl \rangle$$

In the above structure, hl is the hop list and dl is the distance record which consists of distances between two neighbour nodes.

Once RRep is received by the source node ' s ', the RRep is registered into the best paths list by the source node.

3.4. Route Selection

Let's assume we have route discovery path through EER-RBM and discovered three major routes and the shortest route to get to the destination, as shown in Figure 1. However, there is no assurance that the shortest path obtained will be power-efficient path. A pre-determined energy strategy is to quickly disclose an energy-efficient route near to the mainly energy-efficient route and then route it to a preservation plan to increase energy efficiency.

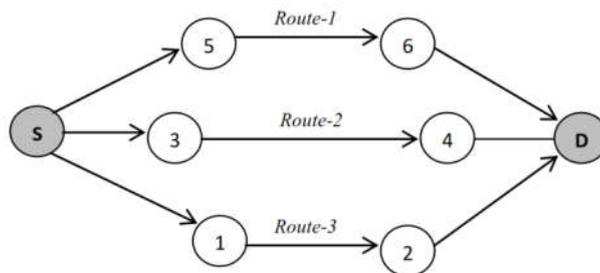


Fig.1 : Route Discovered from S to D by means of EER-RBM

In Fig.1, shows the intermediate hosts between S and D as {1, 2, 3, 4, 5, 6}, and there are 3 shortest routes possible between S and D as, {(S→1→2→D), (S→3→4→D), (S→5→6→D)}.

The d/E is "distance/primary hop energy" defines the optimality condition needed to cover the distance at different power level.

Table- I : Link Cost Table

Next Hop	Distance(d)	Primary Hop Energy(E)	Optimality Criteria(d/E)
S→1	1	9	0.11
S→3	2	9	0.22
S→5	2	9	0.22
1→2	1	5	0.2
3→4	3	7	0.42
5→6	2	6	0.33
2→D	1	8	0.12
4→D	3	9	0.33
6→D	1	6	0.16

Based on the Table we can calculate the possible predetermine energy for each route by adding up all the values obtained in optimality condition from S to D.

Table-II : Predetermine Energy Table

Path	Probable Predetermine Energy Req. (P_E)
S→1→2→D	$P(0.11+0.2+0.12) = 0.43$
S→3→4→D	$P(0.22+0.42+0.33) = 0.97$
S→5→6→D	$P(0.22+0.33+0.16) = 0.71$

Based on the pre-determined energy values in the Table-II, we choose a path which is having minimum value.

4. RESULT EVALUATION

4.1. Experimental Setup

To calculate the enactment we formed a EER-RBM protocol, and using the Table-III simulation parameter over GloMoSim Simulator calculated packet Delivery Ratio and energy efficiency for finishing the broadcasting execution. We also compare the EER-RBM with an existing PEER [16] and DSR to measure Packet Delivery Ratio and Energy Efficiency.

TABLE- III: SIMULATION PARAMETERS

Configuration	Parameter Values
Simulation Area	1000m X 1000m
Packet Size	512 bytes
CBR Rates	4 pkts/sec
Mobility	RWP
Pause time(sec)	60
No. of Nodes	20,40,60,80,100
No. of (S-D) pairs	5,10,20,30,40
Mobility(m/s)	10

4.2. Result Evaluation

With respect to scenario-1, the respective outcomes observed support on the number of node variation are shown in fig.2 between DSR, PEER[16], and EER-RBM.

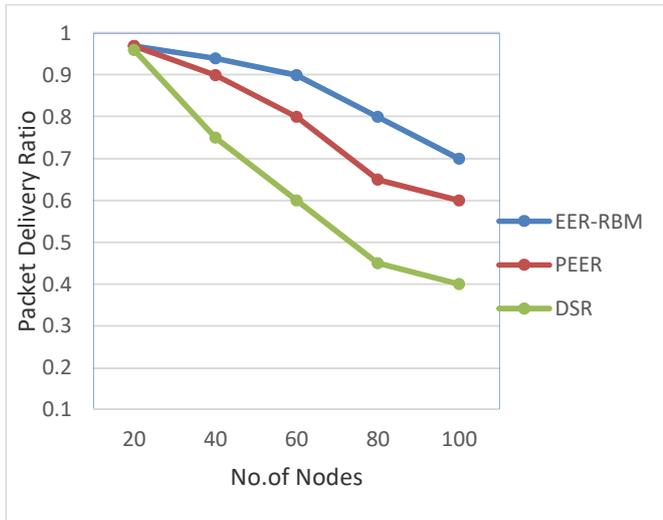


Fig.2. Throughput Comparison

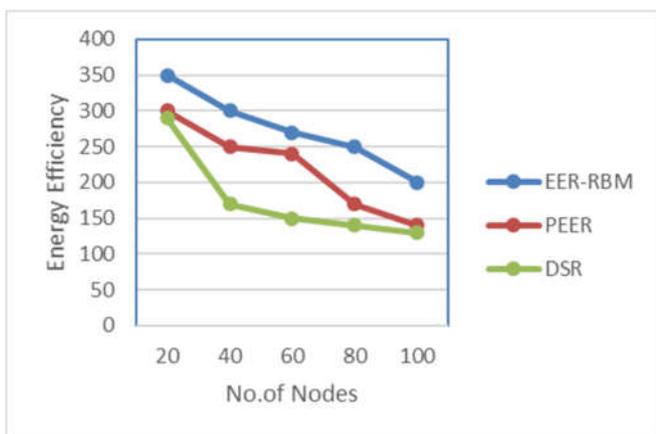


Fig.3. Energy Efficiency Comparison

The comparison results show an inventiveness over traditional DSR and PEER protocol in both the scenarios. In the case of Packet Delivery Ratio all three have to maintain constant fall rate but EER-RBM shows a 25 - 30% better PDR in compare as shown in Fig.2. Fig 3, shows the evaluation similarity of energy utilization efficiency. A high-quality difference in energy efficiency is observed.

5. CONCLUSION

This paper proposes new energy efficient routing model EER-RBM which uses conditional broadcasting and chooses the best path from the paths selected during route discovery phase. The path identified in a predetermined energy efficiency process can be energy efficient enough for efficient data delivery. Experimental results show that energy efficiency is better and has improved throughput. In future, this protocol can further be modified and apply to scalable environment

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